

## THE CHROMATIC NUMBER OF RANDOM GRAPHS

TOMASZ ŁUCZAK

*Received March 9, 1988*

*Revised February 22, 1989*

Let  $\chi(G(n, p))$  denote the chromatic number of the random graph  $G(n, p)$ . We prove that there exists a constant  $d_0$  such that for  $np(n) > d_0$ ,  $p(n) \rightarrow 0$ , the probability that

$$\frac{np}{2 \log np} \left( 1 + \frac{\log \log np - 1}{\log np} \right) < \chi(G(n, p)) < \frac{np}{2 \log np} \left( 1 + \frac{30 \log \log np}{\log np} \right)$$

tends to 1 as  $n \rightarrow \infty$ .

Let  $G(n, p)$  be a random graph with vertex set  $[n] = \{1, 2, \dots, n\}$  in which each possible edge is present independently with the probability  $p = p(n)$ . We say that  $G(n, p)$  has some property a.s. if the probability that it has this property tends to 1 as  $n \rightarrow \infty$ . In this paper we shall consider the asymptotical behaviour of the chromatic number of  $G(n, p)$ .

The question about the value of  $\chi(G(n, p))$  has remained open for over decade until Bollobás in [1] determined  $\chi(G(n, p))$  for  $p(n)$  which either is a constant or tends to 0 slowly enough i.e. for  $p(n) > n^{-1/3+\varepsilon}$  where  $\varepsilon > 0$ . Matula and Kučera gave in [4] an alternative proof for the case when  $p$  is a constant. We shall solve this problem for every function  $p(n) \rightarrow 0$  such that  $np(n)$  is greater than some large constant. Our main result is the following.

**Theorem.** *There exists a constant  $d_0$  such that if  $d = d(n) = np(n) > d_0$  and  $p(n) \rightarrow 0$  then a.s.*

$$\frac{d}{2 \log d} \left( 1 + \frac{\log \log d - 1}{\log d} \right) < \chi(G(n, p)) < \frac{d}{2 \log d} \left( 1 + \frac{30 \log \log d}{\log d} \right).$$

The lower bound is an immediate consequence of the fact that if  $\varepsilon > 0$ , then there exists  $d_\varepsilon$  such that for  $d(n) = np(n) > d_\varepsilon$ , the independence number  $\alpha(G(n, p))$  of  $G(n, p)$  is a.s. less than

$$(1) \quad \alpha(G(n, p)) < \frac{2n}{d} (\log d - \log \log d + 1 - \log 2 + \varepsilon).$$

(This result follows easily from computing the first moment of an appropriate variable.) Thus it is enough to show the second inequality.

Since as we mentioned above the case  $p(n) > n^{-1/3+\varepsilon}$  is studied in [1], we shall show our Theorem only for  $p(n) < \log^{-7} n$ . The method we use works for all  $p(n)$  but this technical assumption will simplify calculations significantly.

**Lemma 1.** Let  $\varepsilon > 0$  and  $k_0 = \lfloor (2n/d) \cdot (\log d - \log \log d + 1 - \log 2 - \varepsilon) \rfloor$ . There exists a constant  $d_\varepsilon$  such that whenever  $n \log^{-7} n > np(n) = d(n) > d_\varepsilon$  then, with the probability at least  $1 - n^{-2}$ ,  $G(n, p)$  contains a subset with at least  $n \log^{-5} d$  vertices which can be properly coloured using  $n \log^{-5} d / k_0$  colours.

**Proof.** Let  $z = \log^{4.9} d$  and  $\mathcal{V} = \mathcal{V}(n, p, z)$  denote the partition of  $[n]$  onto sets  $V_1, \dots, V_{n/z}$  such that  $z - 1 \leq |V_i| \leq z + 1$  for  $i = 1, \dots, n/z$ . We shall call a subset  $S$  of  $G(n, p)$   $\mathcal{V}$ -disjoint if no two elements of  $S$  are contained in the same set of partition  $\mathcal{V}$ .

Define  $\mathcal{G}_i$  as a graph such that subgraphs induced by  $\bigcup_{r=1}^i V_r$  in  $\mathcal{G}_i$  and  $G(n, p)$  are the same and all edges of  $\mathcal{G}_i$  which are not contained in  $\bigcup_{r=1}^i V_r$  are present in this graph independently with the probability  $p$ . Let  $X$  denote the size of the largest  $\mathcal{V}$ -disjoint subset contained in  $G(n, p)$  which can be properly coloured using  $n \log^{-5} d / k_0$  colours, and for  $i = 1, 2, \dots, z$ , let  $X_i$  be the expectation of the size of the largest  $n \log^{-5} d / k_0$  colourable,  $\mathcal{V}$ -disjoint subset in  $\mathcal{G}_i$ . Then  $X = X_{n/z}$ , and setting  $X_0 = EX$  we have

$$|X_i - X_{i-1}| \leq 1 \quad \text{for } i = 1, \dots, n/z.$$

Furthermore the sequence  $X_0, \dots, X_{n/z}$  is a martingale known as *Doob's Martingale Process* (see [5]). Hence, from a martingale inequality of Azuma (see [1] or [5]), we get

$$\text{Prob} \{ |X - EX| \geq t \} \leq 2 \exp \left( -\frac{zt^2}{2n} \right),$$

so

$$\text{Prob} \left\{ |X - EX| \geq \frac{n}{\log^{6.1} d} \right\} \leq 2 \exp \left( -\frac{n}{2 \log^{7.3} d} \right).$$

Thus, to prove Lemma 1, it is enough to show that the probability that  $G(n, p)$  contains a  $n \log^{-5} d / k_0$  colourable,  $\mathcal{V}$ -disjoint subset with more than  $n(\log^{-5} d + \log^{-6.1} d)$  elements is greater than  $\exp(-n \log^{-7.5} d)$  (this observation was first made by Frieze in [2]).

Let  $Y$  be the number of  $n \log^{-5} d / k_0$  colourable,  $\mathcal{V}$ -disjoint subsets of  $mk_0$  elements, where  $m = \lceil n(\log^{-5} d + \log^{-6.1} d) / k_0 \rceil$ , which can be split into exactly  $m$  independent sets, each of  $k_0$  elements. Then

$$P(Y > 0) \geq \frac{(EY)^2}{EY^2}$$

and

$$\begin{aligned}
 \frac{EY^2}{(EY)^2} &\leq \prod_{i=1}^m \sum_{\substack{k_1, k_2, \dots, k_{m+1} \\ \sum_{j=1}^{m+1} k_j = k_0}} \frac{\binom{k_0}{k_1} \binom{k_0}{k_2} \dots \binom{k_0}{k_m} \binom{n/z - (i-1)k_0}{k_{m+1}} z^{k_{m+1}}}{\binom{n/z - (i-1)k_0}{k_0} z^{k_0} (1-p)^{\sum_{j=1}^m \binom{k_j}{2}}} \\
 &\leq \left[ \sum_{l=0}^{k_0} \frac{a_l \binom{n/z - mk_0}{k_0 - l} z^{k_0 - l}}{\binom{n/z - mk_0}{k_0} z^{k_0}} \right]^m \\
 &\leq \left[ \sum_{l=0}^{k_0} \frac{a_l}{(n - (k_0 + 1)mz)^l} \cdot \frac{k_0!}{(k_0 - l)!} \right]^m,
 \end{aligned}$$

where

$$a_l = \sum_{\substack{k_1, k_2, \dots, k_m \\ \sum_{j=1}^m k_j = l}} \binom{k_0}{k_1} \binom{k_0}{k_2} \dots \binom{k_0}{k_m} (1-p)^{-\sum_{j=1}^m \binom{k_j}{2}}.$$

Let  $k_{i_1}, k_{i_2}, \dots, k_{i_r}$  be those from  $k_1, k_2, \dots, k_m$  which are greater than  $2n \log \log d/d$ . Since  $\sum_{j=1}^m k_j = l \leq k_0 < 2n \log d/d$ , so  $r < \log d$ . Thus the number of terms with different sequences  $k_{i_1}, k_{i_2}, \dots, k_{i_r}$  is less than

$$(2) \quad (mk_0)^{\log d} < n^{\log d}.$$

Moreover, for every  $k', k''$  such that  $k' \geq k'' \geq 2n \log \log d/d$  and  $k' + k'' \leq l \leq k_0$ , we have

$$\frac{\binom{k_0}{k'} \binom{k_0}{k''} (1-p)^{-\binom{k'}{2} - \binom{k''}{2}}}{\binom{k_0}{k' + k''} (1-p)^{-\binom{k' + k''}{2}}} < 1.$$

Indeed, when  $2n \log \log d/d \leq k' + k'' \leq 0.7k_0$  then

$$\begin{aligned}
 \frac{\binom{k_0}{k'} \binom{k_0}{k''} (1-p)^{-\binom{k'}{2} - \binom{k''}{2}}}{\binom{k_0}{k' + k''} (1-p)^{-\binom{k' + k''}{2}}} &= \frac{\binom{k_0}{k''} k''}{(k_0 - k') k''} \binom{k' + k''}{k''} \exp(-pk''k') \\
 &\leq \left( \frac{k_0 - k''}{k_0 - k' - k''} \frac{e(k' + k'')}{k''} \exp(-pk') \right)^{k''} \leq \left( \frac{20}{\log d} \right)^{k''} < 1,
 \end{aligned}$$

whereas for  $k' + k'' \geq 0.7k_0$ ,  $k' \geq k'' \geq 2n \log \log d/d$  we have

$$\frac{\binom{k_0}{k'} \binom{k_0}{k''} (1-p)^{-\binom{k'}{2} - \binom{k''}{2}}}{\binom{k_0}{k' + k''} (1-p)^{-\binom{k' + k''}{2}}} \leq 2^{k_0} 2^{k_0} \exp(-pk'k'') < 1.$$

Hence

$$(3) \quad \binom{k_0}{k_{i_1}} \binom{k_0}{k_{i_2}} \dots \binom{k_0}{k_{i_r}} (1-p)^{-\sum_{j=1}^r \binom{k_{i_j}}{2}} \leq \binom{k_0}{l} \exp\left(\frac{l^2 d}{2n}\right).$$

Furthermore, for every choice of  $\bar{k}_1, \bar{k}_2, \dots, \bar{k}_s$ , one can easily get the following inequality

$$(4) \quad \sum_{\substack{\bar{k}_1, \bar{k}_2, \dots, \bar{k}_s \\ \sum_{j=1}^s \bar{k}_j = l \\ \max_{1 \leq j \leq s} \bar{k}_j = f}} \binom{k_0}{\bar{k}_1} \binom{k_0}{\bar{k}_2} \dots \binom{k_0}{\bar{k}_s} (1-p)^{-\sum_{j=1}^s \binom{\bar{k}_j}{2}} \leq \binom{sk_0}{l} \exp\left(\frac{fld}{2n}\right).$$

Thus, from (2), (3) and (4) we obtain

$$a_l \leq \binom{mk_0}{l} \log^{2l} d + n^{\log d} \sum_{i=2n \log \log d/d}^l \binom{k_0}{i} \exp\left(\frac{i^2 d}{2n}\right) \binom{mk_0}{l-i} \log^{2(l-i)} d.$$

For  $2n \log \log d/d \leq i \leq l$ , set

$$b_{i,l} = \binom{k_0}{i} \exp\left(\frac{i^2 d}{2n}\right) \binom{mk_0}{l-i} \log^{2(l-i)} d.$$

Then

$$(5) \quad \begin{aligned} \frac{b_{i+1,l}}{b_{i,l}} &= \frac{k_0 - i}{i+1} \exp\left(\frac{(2i+1)d}{2n}\right) \frac{l-i}{mk_0 - l + i + 1} \frac{1}{\log^2 d} \\ &= \left(\frac{l+1}{i+1} - 1\right) \frac{k_0 - i}{mk_0 - l + i + 1} \exp\left(\frac{(2i+1)d}{2n}\right) \frac{1}{\log^2 d}. \end{aligned}$$

Since the second factor of (5) is less than  $2 \log^6 d/d$  so there exists  $i_1$ , such that  $b_{i+1,l}/b_{i,l} < 1$  for  $i_0 = 2n \log \log d/d < i \leq i_1$ . Then, if  $l$  is large enough, the exponential factor becomes large, and for some  $i_2$   $b_{i+1,l}/b_{i,l} \geq 1$  whenever  $i_1 < i \leq i_2$ . Finally the first factor and the numerator of the second one make  $b_{i+1,l}/b_{i,l}$  less than 1 for  $i_2 < i \leq l$ . Hence,  $b_{i,l}$  is maximised either when  $i = i_0$  or, provided  $i_1 < l$ , for  $i = i_2$ .

An upper bound for  $b_{i_0,l}$  is given by

$$\begin{aligned} b_{i_0,l} &= \binom{k_0}{i_0} \exp\left(\frac{i_0^2 d}{2n}\right) \binom{mk_0}{l-i_0} \log^{2(l-i_0)} d \\ &\leq \binom{(m+1)k_0}{l} \log^{2l} d \leq \left(\frac{3mk_0 \log^2 d}{l}\right)^l. \end{aligned}$$

Let  $l = (\alpha + o(1))k_0$  for some constant  $0 < \alpha \leq 0.9$ . Thus, for  $i' = 0.99l$ , the first factor in (5) tends to a constant, the second is equal  $O(1) \log^6 d/d$ , whereas from the exponential one we get  $d^{1.98\alpha + o(1)}$ . Hence  $b_{i'+1,l}/b_{i',l} < 1$  for  $\alpha < 1/1.98$  and  $i \leq i'$ , whereas, when  $\alpha > 1/1.98$ , we have  $b_{i'+1,l}/b_{i',l} \geq 1$ . Setting  $i' = 0.995l$  one can check in the same way that also for  $\alpha = 1/1.98$   $b_{i,l}$  is maximised for either  $i = i_0$  or  $i_2 > 0.99l$ . Moreover for  $l \leq 0.9k_0$  we have

$$\begin{aligned} \max\{b_{i,l} : 0.99l \leq i \leq l\} &\leq \binom{2k_0}{l} \exp\left(\frac{l^2 d}{2n}\right) \left(\frac{100emk_0}{l}\right)^{0.01l} \log^{2l} d \\ &\leq \left(\frac{2ek_0}{l} \exp\left(\frac{ld}{2n}\right) \left(\frac{n}{l}\right)^{0.01} \log^2 d\right)^l \\ &\leq \left(\frac{6k_0 d^{0.92}}{l}\right)^l. \end{aligned}$$

If  $l \geq 0.9k_0$  then  $i_2 > i'' = l - ld^{-0.3}$ . Indeed, the first factor of  $b_{i''+1,l}/b_{i'',l}$  is bounded above by  $d^{-0.3}$ , the second one is larger than  $d^{-1.3}$  and the exponential factor gives us at least  $d^{1.8}$ . Thus  $b_{i''+1,l}/b_{i'',l} > 1$  and so  $i_2 > i''$ . Furthermore the following inequality holds

$$\begin{aligned} \max\{b_{i,l} : l - ld^{-0.3} \leq i \leq l\} &\leq \binom{k_0}{l - ld^{-0.3}} \exp\left(\frac{l^2 d}{2n}\right) \left(\frac{emk_0}{ld^{-0.3}}\right)^{ld^{-0.3}} \log^{2ld^{-0.3}} d \\ &\leq \binom{k_0}{k_0 - l} \exp\left(\frac{l^2 d}{2n}\right) d^{3k_0 d^{-0.3}}. \end{aligned}$$

Hence

$$\begin{aligned} &\sum_{l=0}^{k_0} \frac{a_l}{(n - (k_0 + 1)mz)^l} \cdot \frac{k_0!}{(k_0 - l)!} \leq \sum_{l=0}^{k_0} \left(\frac{2k_0}{n}\right)^l \left(\frac{3mk_0 \log^2 d}{l}\right)^l \\ &+ n^{2 \log d} \sum_{l=0}^{0.9k_0} \left(\frac{2k_0}{n}\right)^l \left(\frac{6k_0 d^{0.92}}{l}\right)^l \\ &+ n^{2 \log d} d^{3k_0 d^{-0.3}} \sum_{j=0}^{0.1k_0} \frac{k_0!}{(n - (k_0 + 1)mz)^{k_0}} \frac{n^j}{j!} \binom{k_0}{j} \exp\left(\frac{k_0^2 d}{2n}\right) \exp\left(\frac{(-2jk_0 + j^2)d}{2n}\right) \\ &\leq \sum_{l=0}^{k_0} \left(\frac{6mk_0^2 \log^2 d}{ln}\right)^l + n^{2 \log d} \sum_{l=0}^{0.9k_0} \left(\frac{n}{ld^{1.05}}\right)^l \\ &+ n^{3 \log d} d^{3k_0 d^{-0.3}} \left(\frac{k_0}{en(1 - \log^{-0.05} d)} \exp\left(\frac{k_0 d}{2n}\right)\right)^{k_0} \sum_{j=0}^{0.1k_0} \left(\frac{k_0 n}{j^2 d^{1.8}}\right)^j. \end{aligned}$$

All terms of the first sum are less than  $\exp\left(6mk_0^2 \log^2 d/n\right)$ , for  $d$  is large enough, the second sum is bounded above by  $\exp(nd^{-1.04})$ , whereas for the last term we have

$$\frac{k_0}{en(1 - \log^{-0.05} d)} \exp\left(\frac{k_0 d}{2n}\right) < \frac{2(\log d - \log \log d + 1)}{en(1 - \log^{-0.05} d)} \frac{ede^{-\varepsilon}}{2 \log d} < 1$$

and

$$n^{3 \log d} d^{3k_0 d^{-0.3}} \sum_{j=0}^{0.1k_0} \left(\frac{k_0 n}{j^2 d^{1.8}}\right)^j \leq \exp(3k_0 d^{-0.25} + nd^{-1.3}) \leq \exp(nd^{-1.2}).$$

Thus

$$\begin{aligned} \sum_{l=0}^{k_0} \frac{a_l}{(n - (k_0 + 1)mz)^l} \frac{k_0!}{(k_0 - l)!} \\ \leq n^{3 \log d} \exp\left(\frac{6mk_0^2 \log^2 d}{n}\right) + n^{2 \log d} \exp\left(nd^{-1.04}\right) + \exp\left(nd^{-1.2}\right) \\ \leq \exp\left(\frac{7mk_0^2 \log^2 d}{n}\right). \end{aligned}$$

and, finally, we arrive at

$$\begin{aligned} \text{Prob}(Y > 0) &\geq \frac{\mathbb{E} Y^2}{(\mathbb{E} Y)^2} \geq \left[ \exp\left(\frac{7mk_0^2 \log^2 d}{n}\right) \right]^{-m} \\ &\geq \exp\left(-7n \log^2 d \left(\frac{k_0 m}{n}\right)^2\right) > \exp(-n \log^{-7.5} d). \end{aligned}$$

This completes the proof of Lemma 1. ■

**Lemma 2.** *There is a constant  $d_0$  such that for  $d = d(n) = np(n)$ , where  $n \log^{-7} n > d > d_0$ , with the probability greater than  $1 - o(1) - \log^{-1} d$ , more than  $n - 2n \log^{-3} d$  vertices of  $G(n, p)$  can be properly coloured with less than  $\frac{d}{2 \log d} \left(1 + \frac{29 \log \log d}{\log d}\right)$  colours.*

**Proof.** In the proof we shall use an “expose-and-merge” technique introduced by David Matula (for details see [3] and [4]).

For  $A \subset [n]$  define  $[A]^2 = \{\{v, w\} : v, w \in A\}$  and let

$$\bar{k}_0 = \frac{2n}{d}(\log d - 29 \log \log d + 0.1), \quad l_0 = n/(\bar{k}_0 \log^{33} d).$$

Consider the following algorithm:

**Algorithm**

$E := \emptyset$  ;  
 $F_0 := \emptyset$  ;  
 $W_0 := \emptyset$  ;  
**for**  $i = 1$  **to**  $\log^{33} d - \log^{30} d$  **do**  
  **begin**  
    choose randomly  $A_i \subset [n] \setminus W_{i-1}$  with  $|A_i| = n \log^{-28} d$  ;  
    define  $\mathcal{G}_i$  as the graph with the set of vertices  $A_i$  and the set of edges  $E_i$ ,  
      where the probability that  $\{v, w\} \in E_i$  is independent and equal  $p$  for each  
       $\{v, w\} \in [A_i]^2$  ;  
    choose a family  $\{\tilde{\mathcal{R}}_1^i, \tilde{\mathcal{R}}_2^i, \dots, \tilde{\mathcal{R}}_{l_0}^i\}$  of disjoint independent sets from  $A_i$ , such  
      that  $\sum_{l=0}^{l_0} |\tilde{\mathcal{R}}_l^i| = n \log^{-33} d$  – if it is not possible **FAIL** ;  
     $E'_i := E_i \setminus (E_i \cap F_{i-1})$  ;  
     $E := E \cup E'_i$  ;  
     $F_i := F_{i-1} \cup [A_i]^2$  ;  
     $W := W_{i-1} \cup \bigcup_{l=1}^{l_0} \tilde{\mathcal{R}}_l^i$  ;  
  **end**  
 $\bar{F} := [n]^2 \setminus \bigcup_{i=1}^{\log^{33} d - \log^{30} d} F_i$  ;  
  choose  $\bar{E} \subset \bar{F}$  in such a way that each  $e \in \bar{F}$  belong to  $\bar{E}$  with the probability  $p$   
  independently of each other ;  
   $E := E \cup \bar{E}$  ;  
**output**  $E$  ;  $\{\tilde{\mathcal{R}}_1^1, \tilde{\mathcal{R}}_2^1, \dots, \tilde{\mathcal{R}}_{l_0}^{\log^{33} d - \log^{30} d}\}$  ;  
**end**

Let us observe first that the probability that  $\{v, w\} \in E$  is equal  $p$  independently for each  $\{v, w\} \in [n]^2$ , so the graph  $\tilde{\mathcal{G}}$  with the set of vertices  $[n]$  and the set of edges  $E$  may be treated as  $G(n, p)$ .

Obviously, we may consider each  $\mathcal{G}_i$  as  $G(\bar{n}, p)$ , where  $\bar{n} = n \log^{-28} d$ . The average degree of such a random graph is given by  $\bar{d} = p \bar{n} \log^{-28} d = d \log^{-28} d$  and

$$\bar{k}_0 = \frac{2n}{d} (\log d - 29 \log \log d + 0.1) < \frac{2\bar{n}}{\bar{d}} (\log \bar{d} - \log \log \bar{d} + 0.2).$$

Thus, from Lemma 1, the probability that  $\mathcal{G}_i$  contains no subsets with  $n \log^{-33} d < \bar{n} \log^{-5} \bar{d}$  elements which can be properly coloured using  $\bar{n} \log^{-5} \bar{d} / \bar{k}_0$  colours is less than  $n^{-2}$ , so the probability of **FAIL** in the Algorithm is less than  $n^{-1}$ .

Thus, with the probability at least  $1 - o(1)$ , the Algorithm finds

$$(\log^{33} d - \log^{30} d) l_0 = \frac{n - n \log^{-3} d}{\bar{k}_0} < \frac{d}{2 \log d} \left( 1 + \frac{29 \log \log d}{\log d} \right)$$

disjoint sets  $\tilde{\mathcal{R}}_1^i, \tilde{\mathcal{R}}_2^i, \dots, \tilde{\mathcal{R}}_{l_0}^i$  such that

$$\sum_i \sum_l |\tilde{\mathcal{R}}_l^i| = (\log^{33} d - \log^{30} d) n \log^{-33} d = n - n \log^{-3} d.$$

Note however, that although  $\tilde{\mathcal{R}}_l^i$  is an independent subset of  $\mathcal{G}_i$  it is *not* necessarily independent as a subset of  $\tilde{\mathcal{G}}$ . Let  $X$  denote the number of edges of  $\tilde{\mathcal{G}}$  contained in  $\tilde{\mathcal{R}}_l^i$  for some  $1 \leq i \leq \log^{33} d - \log^{30} d$ ,  $1 \leq l \leq l_0$ . We shall estimate from above the size of  $X$ .

Let  $\{v, w\} \in [n]^2$  be such that  $\{v, w\} \in E$  and  $\{v, w\} \in \tilde{\mathcal{R}}_l^i$  for some  $i, l$ . Denote by  $i(v, w)$  the smallest number  $i$  for which  $\{v, w\} \subset A_i$  and let  $j(v, w), l(v, w)$  be such that  $\{v, w\} \subset \tilde{\mathcal{R}}_{l(v, w)}^{j(v, w)}$ . Notice that  $i(v, w) < j(v, w)$  since  $\{v, w\} \in E$  implies that  $\{v, w\}$  is an edge of  $\mathcal{G}_{i(v, w)}$  whereas the set  $\tilde{\mathcal{R}}_{l(v, w)}^{j(v, w)}$  contains no edges of  $\mathcal{G}_{j(v, w)}$ . Since for all  $i$  we have  $|W_i| \leq n - n \log^{-3} d$ , so the probability that for chosen  $i(v, w), j(v, w)$ , a pair  $\{v, w\}$  is contained in both  $A_{i(v, w)}$  and  $A_{j(v, w)}$  is less than  $(2 \log^{-25} d)^4$ . Now observe that for any  $l$ , where  $1 \leq l \leq l_0$ , each subset of  $A_{j(v, w)}$  of  $|\tilde{\mathcal{R}}_l^{j(v, w)}|$  elements is equally likely to be chosen as  $\tilde{\mathcal{R}}_l^{j(v, w)}$  (i.e. this event depends only on the structure of  $\mathcal{G}_{j(v, w)}$  which is symmetric with respect to the labelling of vertices). Moreover, due to (1), we may assume that for all  $i, l$ , we have  $|\tilde{\mathcal{R}}_l^i| < 2n \log d/d$ . Thus, since  $|A_{j(v, w)}| = n \log^{-28} d$ , the probability that both  $v, w$  are in the same set  $\tilde{\mathcal{R}}_{l(v, w)}^{j(v, w)}$  for some  $l(v, w)$  is less than  $2 \log^{29} d/d$ .

So finally we arrive at the following upper bound for the expectation of  $X$

$$\mathbb{E} X \leq \binom{n}{2} \binom{\log^{33} d - \log^{30} d}{2} \cdot 16 \log^{-100} d \cdot p \cdot 2 \log^{29} d/d < 8n \log^{-5} d.$$

Thus, from Markov's inequality,

$$\text{Prob}(X > 0.5n \log^{-3} d) < \log^{-1} d.$$

Now, for all  $i, l$ , delete from  $\tilde{\mathcal{R}}_l^i$  all these vertices which belongs to edges of  $\tilde{\mathcal{G}}$  which are contained in  $\tilde{\mathcal{R}}_l^i$  and denote the set obtained in this way by  $\mathcal{R}_l^i$ . Then

$$\sum_{i, l} |\mathcal{R}_l^i| \geq \sum_{i, l} |\tilde{\mathcal{R}}_l^i| - 2X = n - n \log^{-3} d - 2X$$

so

$$\text{Prob} \left( \sum_{i, l} |\mathcal{R}_l^i| \geq n - 2n \log^{-3} d \right) > 1 - o(1) - \log^{-1} d$$

and the assertion follows. ■



**Proof of Theorem.** Lemma 2 implies that with the probability at least  $1 - o(1) - \log^{-1} d$  we can colour  $n - 2n \log^{-3} d$  vertices of  $G(n, p)$  using only  $\frac{d}{2 \log d} \left(1 + \frac{29 \log \log d}{\log d}\right)$  colours. One can easily check that for  $d$  large enough a.s. every subgraph of  $G(n, p)$  on  $s$  vertices, where  $s \leq 2n \log^{-3} d$ , has less than  $sd \log^{-2.5} d$  edges. Hence, for large  $d$ , we have

$$(6) \quad \text{Prob} \left( \chi(G(n, p)) \leq \frac{d}{2 \log d} \left(1 + \frac{29.5 \log \log d}{\log d}\right) \right) > 1 - o(1) - \log^{-1} d,$$

and for  $d(n) \rightarrow \infty$ , the assertion follows.

Moreover, Shamir and Spencer proved in [5] that for every  $d(n) < \log n$  there exists a function  $u_d(n)$  such that a.s.

$$u_d(n) \leq \chi(G(n, p)) \leq u_d(n) + 4.$$

Thus, since from (6) for  $d(n)$  greater than some constant  $d_0$  we have

$$\liminf_{n \rightarrow \infty} \text{Prob} \left( \chi(G(n, p)) \leq \frac{d}{2 \log d} \left(1 + \frac{29.5 \log \log d}{\log d}\right) \right) > 0.5$$

so, for such  $d(n)$ , a.s.

$$\chi(G(n, p)) \leq \frac{d}{2 \log d} \left(1 + \frac{29.5 \log \log d}{\log d}\right) + 5 \leq \frac{d}{2 \log d} \left(1 + \frac{30 \log \log d}{\log d}\right).$$

This completes the proof of Theorem. ■

**Acknowledgement.** I wish to thank Alan Frieze and Joel Spencer for their helpful comments and suggestions.

## References

- [1] B. BOLLOBÁS: The chromatic number of random graphs, *Combinatorica*, **8** (1988), 49–56.
- [2] A. M. FRIEZE: On the independence number of random graphs, *Disc. Math.*, **81** (1990), 171–175.
- [3] D. MATULA: Expose-and-merge exploration and the chromatic number of a random graph, *Combinatorica*, **7** (1987), 275–284.
- [4] D. MATULA, and L. KUČERA: An expose-and-merge algorithm and the chromatic number of a random graph, in “Proceedings of Random Graphs ’87”, Wiley, Chichester, **1990**, 175–188.

- [5] E. SHAMIR, and J. SPENCER: Sharp concentration of the chromatic number on random graphs  $G_{n,p}$ , *Combinatorica*, **7** (1987), 124–129

**Tomasz Łuczak**

*Institute for Mathematics and its Applications,  
University of Minnesota,  
Minneapolis,  
U.S.A.*

*Permanent address:  
Institute of Mathematics,  
Adam Mickiewicz University,  
Poznań,  
Poland*